

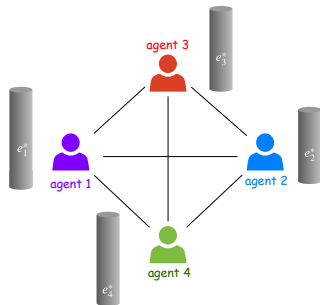
Incentivizing Effort in Interdependent Security Games Using Resource Pooling

Mohammad Mahdi Khalili, Xueru Zhang, Mingyan Liu

University of Michigan, Ann Arbor

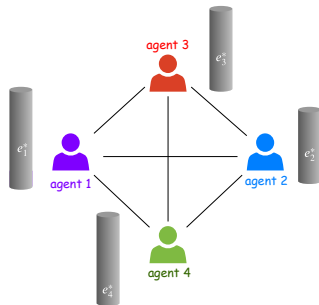
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InterDependent Security (IDS) Games



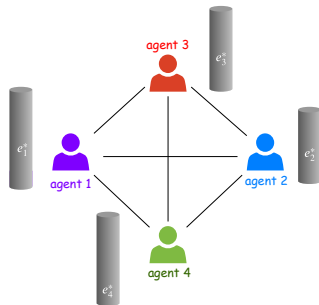
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- Players: Agents/firms



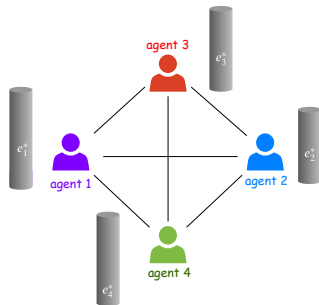
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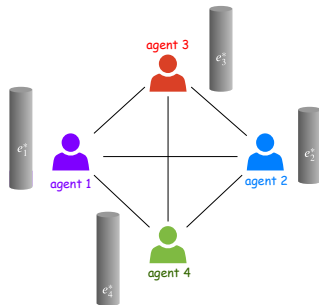
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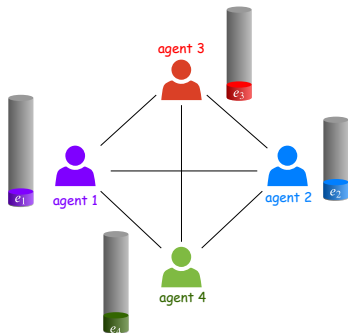
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- Players: Agents/firms
- Actions: Security investment/effort e_i
- Utility: Positive externality
- Socially optimal effort level: e_i^*
- Under-investment: $e_i < e_i^*$



Under-investment

Reason of under-investment in security

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- Taking advantage of other firms effort and investment

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Examples of under-investment

- Equifax: was aware of software vulnerability 2 months before the data breach
- JPMorgan Chase: one of the servers did not have 2 factors authentication

Addressing Under-investment Issue

Traditional Solutions in the literature of IDS games:

Addressing Under-investment Issue

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- Incentive/Taxation Mechanism ¹

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Traditional Solutions in the literature of IDS games:

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 - Costly due to the social/privacy reasons

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IDS game with resource pooling (A non-cooperative game):

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Examples of Resource pooling,

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- Security product discount

Overview

Motivation

IDS Game Without Resource Pooling

IDS Game with Resource Pooling

Voluntary Participation

Community Based Resource Pooling

Conclusion

Model: IDS Game without Resource Pooling

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- Agent i 's utility

$$u_i(e_i, e_{-i}) = -l_i + a_i e_i - b_i e_i^2 + e_i \sum_j x_{ij} e_j$$

No Resource Pooling: Nash Equilibrium

Theorem

Assume $2b_i \geq \sum_j x_{ij}, \forall i$. Then, IDS game without resource pooling has the unique NE ($\hat{\mathbf{e}} = [\hat{e}_1, \hat{e}_2, \dots, \hat{e}_n]$).

No Resource Pooling: Socially Optimal Effort Level

- Socially Optimal Effort level

$$\mathbf{e}^* = \arg \max_{\mathbf{e}=[e_1, e_2, \dots, e_n]} \sum_i u_i(\mathbf{e})$$

Theorem

Assume $2b_i \geq \sum_j x_{ij} + x_{ji}, \forall i$. Then, socially optimal effort level \mathbf{e}^ is unique. Moreover, $e_i^* > \hat{e}_i, \forall i$.*

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$$v_i(\mathbf{e}_i, \mathbf{e}_{-i}) = -l_i + a_i E_i + E_i \cdot \sum_{j=1}^n x_{ij} E_j - \sum_{k=1}^n b_k \cdot e_{ik}^2.$$

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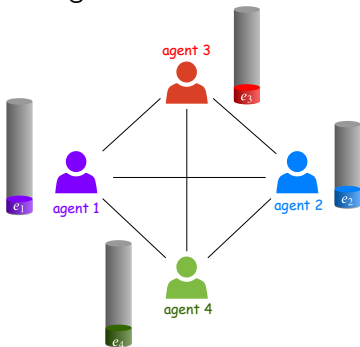
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Let \hat{E}_i be the total effort exerted on behalf of agent i at the NE.

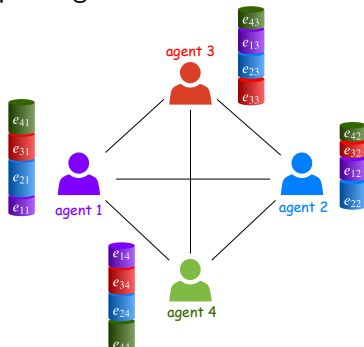
Then, $\begin{bmatrix} \hat{E}_1 \\ \vdots \\ \hat{E}_n \end{bmatrix} = \mathbf{e}^*$ is the total effort at the NE of the game with resource pooling.

Resource Pooling: Nash Equilibrium

IDS game **without** resource pooling



IDS game **with** resource pooling



Resource Pooling: Nash Equilibrium

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Let $2b_i > \sum_j x_{ji} + x_{ij}, \forall i$ and $\hat{E} = [\hat{e}_{ij}]$ be the effort profile at the NE of the game with resource pooling. Then, we have

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- $v_i(\hat{E}) \geq u_i(\hat{e}_i, \hat{e}_{-i})$
- $\sum_{i=1}^n v_i(\hat{E}) \geq \sum_{i=1}^n u_i(\mathbf{e}^*)$

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- Let $\bar{E} = [\bar{e}_{ij}]_{n \times n}$ be the NE of game G^k and $v_k(\bar{E})$ be the utility of agent i at the NE.
- We say that resource pooling has the voluntary participation property with respect to agent k , if

$$v_k(\bar{E}) \leq v_k(\hat{E}), \quad (1)$$

where \hat{E} is the effort profile at the NE of game with resource pooling.

Voluntary Participation

Theorem

Resource Pooling always satisfies the Voluntary Participation with respect to all agents.

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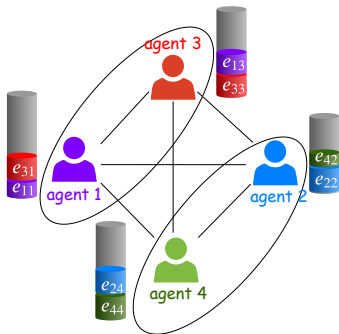
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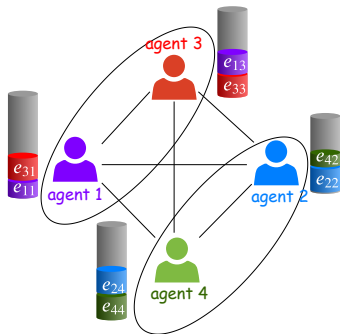
Community Based Resource Pooling

Agents form communities C_1, C_2, \dots, C_m and are allowed to pool resources within their own communities. $C_i \cap C_j = \emptyset$.



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Theorem (Informal)

Community Based Resource Pooling improves agents' utilities and their efforts.

Resource Pooling within Communities: example

$n = 10$ agents in the network

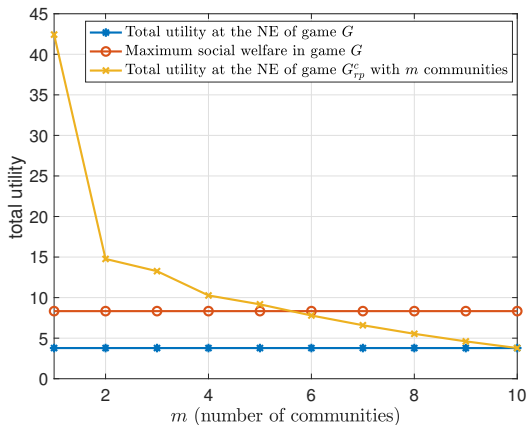


Figure: Total utility as a function of number of communities

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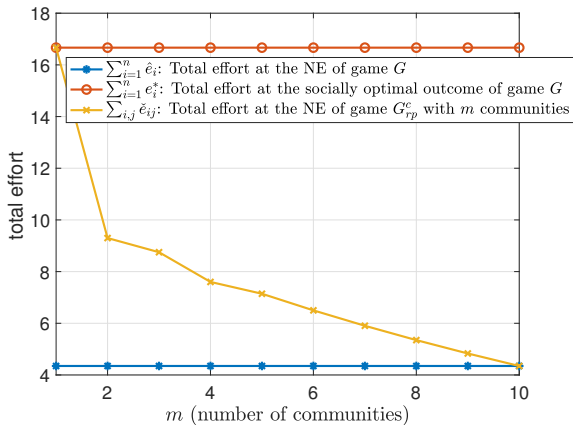


Figure: Total effort as a function of number of communities.

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- resource pooling increases the total effort exerted on behalf of each agent as compared to no resource pooling
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- Social welfare at the NE of the game with resource pooling is higher than the optimal social welfare under the game without resource pooling
- agents voluntarily participate in resource pooling

Future Work

- Study resource pooling with non-quadratic
- Resource pooling may not help under other models
 - With the limited effort budget, resource pooling may not help³

³Khalili, Mohammad Mahdi, Xueru Zhang, and Mingyan Liu. "Public Good Provision Games on Networks with Resource Pooling." *Network Games, Control, and Optimization*. Birkhuser, Cham, 2019. 271-287.   

Mohammad Mahdi Khalili

Email Address

khalili@umich.edu

Home Page

<http://www-personal.umich.edu/~khalili/>